## **Amendments to the Claims**

The following listing of the claims will replace all prior versions, and listings of the claims in the application:

## **Listing of Claims**

## 1-39 Canceled

- 40. (Currently amended) A method for configuring an equivalent  $2^n \times 2^n$  k-stage bit-permuting network based on a given  $2^n \times 2^n$  k-stage bit-permuting network having the  $\underline{a}$  representation  $[\sigma_0:\sigma_1:\sigma_2:...:\sigma_{k-1}:\sigma_k]_n$ , the method comprising: specifying a permutation  $\kappa$  on integers from 1 to n that preserves n, and implementing the equivalent network as  $[\sigma_0:\sigma_1:...:\sigma_{j-1}\kappa:\kappa^{-1}\sigma_j:...:\sigma_k]_n$ , j=1,2,..., or k.
- 41. (Previously presented) The method as recited in claim 40 wherein the given network is a banyan-type network and the equivalent network is a banyan-type network.
- 42. (Currently amended) A method for configuring an equivalent  $2^n \times 2^n$  k-stage bit-permuting network based on a given  $2^n \times 2^n$  k-stage bit-permuting network having the  $\underline{a}$  representation  $[\sigma_0 : \sigma_1 : \sigma_2 : \ldots : \sigma_{k-1} : \sigma_k]_n$ , the method comprising:

specifying permutations  $\kappa_1, \, \kappa_2, \, \ldots \,$ ,  $\kappa_k$  on integers from 1 to n that preserve n, and implementing the equivalent network as  $[\sigma_0 \kappa_1 : \kappa_1^{-1} \sigma_1 \kappa_2 : \kappa_2^{-1} \sigma_2 \kappa_3 : \ldots : \kappa_{k-1}^{-1} \sigma_{k-1} \kappa_k : \kappa_k^{-1} \sigma_k]_n$ .

43. (Previously presented) The method as recited in claim 42 wherein the given network is a banyan-type network and the equivalent network is a banyan-type network.

44. (Currently amended) A method for configuring an equivalent  $2^n \times 2^n$  bit-permuting network based on a given  $2^n \times 2^n$  bit-permuting network composed of stages and exchanges, the method comprising:

identifying one stage from the stages of the given network, the identified stage having a preceding exchange immediately before it and a succeeding exchange immediately after it,

specifying a permutation on the integers 1 to n that preserves n,

rearranging the preceding exchange and the succeeding exchange with reference to the permutation to generate a rearranged preceding exchange and a rearranged succeeding exchange, respectively, and

implementing the equivalent network so that a stage in the equivalent network corresponding to the identified stage has the rearranged preceding exchange and the rearranged succeeding exchange.

- 45. (Previously presented) The method as recited in claim 44 wherein the permutation, denoted as  $\kappa$ , induces a  $2^n \times 2^n$  cell rearrangement  $X_{\kappa}$ , and the rearranging includes multiplying the preceding exchange by  $X_{\kappa}$  from the right-hand side to produce the rearranged preceding exchange and multiplying the succeeding exchange by  $X_{\kappa^{-1}}$  from the left-hand side to produce the rearranged succeeding exchange.
- 46. (Previously presented) The method as recited in claim 45 wherein the given network has k-stages, the given network has the representation  $[\sigma_0:\sigma_1:\sigma_2:\ldots:\sigma_{k-1}:\sigma_k]_n$ , the identified stage is stage j, and the equivalent network is of the form,  $[\sigma_0:\sigma_1:\ldots:\sigma_{j-1}\kappa:\kappa^{-1}\sigma_j:\ldots:\sigma_k]_n$ ,  $j=1,2,\ldots$ , or k.
- 47. (Previously presented) The method as recited in claim 44 wherein the given network is a banyan-type network and the equivalent network is a banyan-type network.

48. (Currently amended) A method for configuring an equivalent  $2^n \times 2^n$  bit-permuting network by cell rearrangement based on a given  $2^n \times 2^n$  bit-permuting network composed of stages and exchanges, the method comprising:

identifying one stage from the stages of the given network, the identified stage having a preceding exchange and a succeeding exchange,

specifying a permutation, denoted as  $\kappa$ , on the integers 1 to n that preserves n and induces a  $2^n \times 2^n$  cell rearrangement  $X_{\kappa}$ ,

rearranging the preceding exchange of the identified stage by multiplying the preceding exchange with  $X_{\kappa}$  from the right-hand side to produce a rearranged preceding exchange and rearranging the succeeding exchange of the identified stage by multiplying the succeeding exchange by  $X_{\kappa^{-1}}$  from the left-hand side to produce a rearranged succeeding exchange, and

implementing the equivalent network so that a stage in the equivalent network corresponding to the identified stage has the rearranged preceding exchange and the rearranged succeeding exchange.

49. (Currently amended) A method for cell rearrangement of a 2<sup>n</sup>×2<sup>n</sup> bit-permuting network composed of stages and exchanges, the method comprising:

selecting one stage from the stages of the given network to identify a preceding exchange and a succeeding exchange,

specifying a permutation, denoted as  $\kappa$ , on the integers 1 to n that preserves n and induces a  $2^n \times 2^n$  cell rearrangement  $X_{\kappa}$ , and

multiplying the preceding exchange with  $X_{\kappa}$  from the right-hand side to implement a rearranged preceding exchange and multiplying the succeeding exchange by  $X_{\kappa^{-1}}$  from the left-hand side to implement a rearranged succeeding exchange.

50. (Currently amended) A method for cell rearrangement of a given stage of a 2<sup>n</sup>×2<sup>n</sup>

bit-permuting network composed of stages and exchanges, the method comprising:

specifying a permutation, denoted as  $\kappa$ , on the integers 1 to n that preserves n and induces a  $2^n \times 2^n$  cell rearrangement  $X_{\kappa}$ , and

multiplying the  $\underline{a}$  preceding exchange immediately before the given stage by  $X_{\kappa}$  from the right-hand side to implement a rearranged preceding exchange for the given stage and multiplying the  $\underline{a}$  succeeding exchange immediately after the given stage exchange by  $X_{\kappa^{-1}}$  from the left-hand side to implement a rearranged succeeding exchange for the given stage.

51. (Currently amended) A method for rearranging a given  $2^n \times 2^n$  k-stage bit-permuting network having the <u>a</u> representation  $[\sigma_0 : \sigma_1 : \sigma_2 : \ldots : \sigma_{k-1} : \sigma_k]_n$  to an equivalent  $2^n \times 2^n$  bit-permuting network having the representation  $[\pi_0 : \pi_1 : \pi_2 : \ldots : \pi_{k-1} : \pi_k]_n$ , the method comprising:

determining permutations  $\kappa_1, \, \kappa_2, \, \ldots, \, \kappa_k$  on integers from 1 to n that preserve n, and implementing the equivalent network with exchanges determined from  $\pi_1 = \kappa_1^{-1} \sigma_1 \kappa_2, \, \pi_2 = \kappa_2^{-1} \sigma_2 \kappa_3, \, \ldots, \, \pi_{k-1} = \kappa_{k-1}^{-1} \sigma_{k-1} \kappa_k \text{ so that the equivalent network can be further expressed as } [\alpha: \kappa_1^{-1} \sigma_1 \kappa_2: \kappa_2^{-1} \sigma_2 \kappa_3: \ldots: \kappa_{k-1}^{-1} \sigma_{k-1} \kappa_k: \beta]_n \text{ for some permutations } \alpha \text{ and } \beta.$ 

- 52. (Previously presented) The method as recited in claim 51 wherein the input exchange  $\alpha$  of the equivalent network is equal to  $\pi_0$ .
- 53. (Previously presented) The method as recited in claim 51 wherein the output exchange  $\beta$  of the equivalent network is equal to  $\pi_k$ .
- 54. (Previously presented) The method as recited in claim 51 wherein the input exchange  $\alpha$  of the equivalent network is equal to  $\pi_0$  and the output exchange  $\beta$  of the equivalent network is equal to  $\pi_k$ .
  - 55. (Currently amended) A method for configuring a given 2<sup>n</sup>×2<sup>n</sup> k-stage bit-permuting

network to achieve a desired trace, the method comprising:

determining a permutation  $\sigma$  on the integers 1 to n that maps the <u>a</u> trace of the given network term-by-term to the desired trace, and

prepending the given network with an extra input exchange induced by  $\sigma^{-1}$  if the permutation  $\sigma$  exists.

- 56. (Previously presented) A method as recited in claim 55 wherein k = n and the bit-permuting network is a  $2^n \times 2^n$  banyan-type network.
- 57. (Previously presented) A method as recited in claim 55 wherein the trace of the given network is the sequence  $t_1, t_2, ..., t_k$ , the desired trace is the sequence  $t'_1, t'_2, ..., t'_k$ , and  $t'_j = \sigma(t_j)$  for j = 1, 2, ..., k.
- 58. (Currently amended) A method for configuring a given  $2^n \times 2^n$  k-stage bit-permuting network to achieve a desired guide, the method comprising:

determining a permutation  $\pi$  on the integers 1 to n that maps the <u>a</u> guide of the given network term-by-term to the desired guide, and

appending the given network with an extra output exchange induced by  $\pi$  if the permutation  $\pi$  exists.

- 59. (Previously presented) A method as recited in claim 58 wherein k = n and the bit-permuting network is a  $2^n \times 2^n$  banyan-type network.
- 60. (Previously presented) A method as recited in claim 58 wherein the guide of the given network is the sequence  $g_1, g_2, ..., g_k$ , the desired guide is the sequence  $g'_1, g'_2, ..., g'_k$ , and  $g'_j = \pi(g_j)$  for j = 1, 2, ..., k.

61. (Currently amended) A method for configuring a given  $2^n \times 2^n$  k-stage bit-permuting network to achieve a desired trace and a desired guide, the method comprising:

determining a permutation  $\sigma$  on the integers 1 to n that maps the  $\underline{a}$  trace of the given network term-by-term to the desired trace,

determining a permutation  $\pi$  on the integers 1 to n that maps the <u>a</u> guide of the given network term-by-term to the desired guide, and

if both the permutations  $\sigma$  and  $\pi$  exist, prepending the given network with an extra input exchange induced by  $\sigma^{-1}$ , and appending the given network with an extra output exchange induced by  $\pi$ .

- 62. (Previously presented) A method as recited in claim 61 wherein k = n and the bit-permuting network is a  $2^n \times 2^n$  banyan-type network.
- 63. (Previously presented) A method as recited in claim 61 wherein the trace of the given network is the sequence  $t_1, t_2, ..., t_k$ , the desired trace is the sequence  $t'_1, t'_2, ..., t'_k$ , and  $t'_j = \sigma(t_j)$  for j = 1, 2, ..., k and wherein the guide of the given network is the sequence  $g_1, g_2, ..., g_k$ , the desired guide is the sequence  $g'_1, g'_2, ..., g'_k$ , and  $g'_j = \pi(g_j)$  for j = 1, 2, ..., k.
- 64. (Currently amended) A method for rearranging a given  $2^n \times 2^n$  banyan-type network having the <u>a</u> representation  $[\sigma_0 : \sigma_1 : \sigma_2 : ... : \sigma_{n-1} : \sigma_n]_n$  to an equivalent  $2^n \times 2^n$  banyan-type network having the representation  $[\pi_0 : \pi_1 : \pi_2 : ... : \pi_{n-1} : \pi_n]_n$ , the method comprising:

determining permutations  $\kappa_1, \kappa_2, \ldots, \kappa_n$  on integers from 1 to n that preserve n, and implementing the equivalent network with exchanges determined from  $\pi_1 = \kappa_1^{-1} \sigma_1 \kappa_2$ ,  $\pi_2 = \kappa_2^{-1} \sigma_2 \kappa_3, \ldots, \pi_{n-1} = \kappa_{n-1}^{-1} \sigma_{n-1} \kappa_n$  so that the equivalent network can be further expressed as  $[\alpha : \kappa_1^{-1} \sigma_1 \kappa_2 : \kappa_2^{-1} \sigma_2 \kappa_3 : \ldots : \kappa_{n-1}^{-1} \sigma_{n-1} \kappa_n : \beta]_n$  for some permutations  $\alpha$  and  $\beta$ .

65. (Previously presented) The method as recited in claim 64 wherein the input

exchange  $\alpha$  of the equivalent network is equal to  $\pi_0$ .

- 66. (Previously presented) The method as recited in claim 64 wherein the output exchange  $\beta$  of the equivalent network is equal to  $\pi_k$ .
- 67. (Previously presented) The method as recited in claim 64 wherein the input exchange  $\alpha$  of the equivalent network is equal to  $\pi_0$  and the output exchange  $\beta$  of the equivalent network is equal to  $\pi_k$ .
- 68. (Currently amended) A method for rearranging a first  $2^n \times 2^n$  banyan-type network having the  $\underline{a}$  representation  $[\sigma_0 : \sigma_1 : \ldots : \sigma_{n-1} : \sigma_n]$  with a first trace induced by a permutation  $\tau$  on integers 1 to n and a first guide induced by a permutation  $\gamma$  on integers 1 to n to a second  $2^n \times 2^n$  banyan-type network having the representation  $[\lambda \sigma_0 : \sigma_1 : \ldots : \sigma_{n-1} : \sigma_n \pi]$ , the method comprising:

prepending an additional input exchange  $X_{\lambda}$  to the first network, and

appending an additional output exchange  $X_{\pi}$  to the first network, wherein the second network is characterized by a second trace induced by a permutation  $\tau'$  on integers 1 to n and a second guide induced by a permutation  $\gamma'$  on integers 1 to n such that  $\tau' = \tau \lambda^{-1}$  and  $\gamma' = \gamma \pi$ .

- 69. (Previously presented) The method as recited in claim 68 wherein the permutations  $\tau$  and  $\gamma$  that induce the first trace and the first guide are converted to any  $\tau'$  and  $\gamma'$ , respectively, with the prepended input exchange  $X_{\lambda}$  and the appended output exchange  $X_{\pi}$  by computing  $\lambda = \tau'^{-1}\tau$  and  $\pi = \gamma^{-1}\gamma'$ .
- 70. (Currently amended) A method for configuring a given  $2^n \times 2^n$  banyan-type network to achieve a desired trace wherein the <u>a</u> trace of the given network is induced by a permutation  $\tau$  on integers 1 to n, and the desired trace is induced by another permutation  $\tau'$  on integers 1 to n,

the method comprising:

determining a permutation  $\lambda = \tau^{-1}\tau$ , and prepending the given network with an extra input exchange induced by  $\lambda$ .

- 71. (Previously presented) A method as recited in claim 70 wherein the desired trace is 1, 2, ..., n and the permutation  $\lambda = \tau$ .
- 72. (Previously presented) A method as recited in claim 70 wherein the desired trace is n, n-1, ..., 1 and the permutation  $\lambda = \sigma_{\leftrightarrow}^{(n)} \tau$ .
- 73. (Currently amended) A method for configuring a given  $2^n \times 2^n$  banyan-type network to achieve a desired guide wherein the <u>a</u> guide of the given network is induced by a permutation  $\gamma$  on integers 1 to n, and the desired guide is induced by another permutation  $\gamma'$  on integers 1 to n, the method comprising:

determining a permutation  $\pi = \gamma^{-1} \gamma'$ , and appending the given network with an extra output exchange induced by  $\pi$ .

- 74. (Previously presented) A method as recited in claim 73 wherein the desired guide is 1, 2, ..., n and the permutation  $\pi = \gamma^{-1}$ .
- 75. (Previously presented) A method as recited in claim 73 wherein the desired guide is n, n-1, ..., 1 and the permutation  $\pi = \gamma^{-1} \sigma_{\leftrightarrow}^{(n)}$ .
- 76. (Currently amended) A method for configuring a given  $2^n \times 2^n$  banyan-type network to achieve a desired trace and a desired guide wherein the <u>a</u> trace of the given network is induced by a permutation  $\tau$  on integers 1 to n, the desired trace is induced by another permutation  $\tau'$  on

integers 1 to n, the <u>a</u> guide of the given network is induced by a permutation  $\gamma$  on integers 1 to n, and the desired guide is induced by another permutation  $\gamma'$  on integers 1 to n, the method comprising:

determining a permutation  $\lambda = \tau'^{-1}\tau$ , determining a permutation  $\pi = \gamma^{-1} \gamma'$ , prepending the given network with an extra input exchange induced by  $\lambda$ , and appending the given network with an extra output exchange induced by  $\pi$ .

77. (Currently amended) An equivalent  $2^n \times 2^n$  k-stage bit-permuting network based on a given  $2^n \times 2^n$  k-stage bit-permuting network having the  $\underline{a}$  representation  $[\sigma_0 : \sigma_1 : \sigma_2 : \ldots : \sigma_{k-1} : \sigma_k]_n$  the equivalent network comprising:

permutation means for computing a permutation  $\kappa$  on integers from 1 to n that preserves n, and

a  $2^n \times 2^n$  k-stage bit-permuting network configured as  $[\sigma_0 : \sigma_1 : \ldots : \sigma_{j-1} \kappa : \kappa^{-1} \sigma_j : \ldots : \sigma_k]_n$ ,  $j = 1, 2, \ldots$ , or k.

78. (Currently amended) An equivalent  $2^n \times 2^n$  k-stage bit-permuting network based on the j-th stage of a given  $2^n \times 2^n$  k-stage bit-permuting network having the <u>a</u> representation  $[\sigma_0 : \sigma_1 : \sigma_2 : ... : \sigma_{k-1} : \sigma_k]_n$  and based on a permutation  $\kappa$  on integers from 1 to n that preserves n, the equivalent network comprising:

an input exchange  $\sigma_0 \kappa$  if j=l, or an input exchange  $\sigma_0$  if j=2,3,...,k, an output exchange  $\kappa^{-1} \sigma_k$  if j=k, or an output exchange  $\sigma_k$  if j=1,2,...,k-l, and interstage exchanges  $\sigma_1,\sigma_2,...,\sigma_{j-l}\kappa,\kappa^{-1}\sigma_j,...,\sigma_{k-l}$  if j=2,..., or k-1, or interstage exchanges  $\kappa^{-1}\sigma_1,\sigma_2,...,\sigma_j,...,\sigma_{k-1}$  if j=1, or interstage exchanges  $\sigma_1,\sigma_2,...,\sigma_j,...,\sigma_{k-2},\sigma_{k-1}\kappa$  if j=k.